## **M.A./M.Sc. Semester III Examination, 2019 (Old pattern under CDOE) Subject: Mathematics (Applied Stream)**

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Time: 2 Hours Full Marks: 45

*The figures in the margin indicate full marks. Candidates are required to give their answers in their own words as far as practicable. [Notation and symbols have their usual meaning]*

### **Paper: MAS305 (Advanced Operations Research-I)**

Answer any five questions. Only first five answers will be evaluated. 9×5 = 45  
\n1. (a) State and prove the sufficient condition for optimality of Lagrange multiplier method for [4]  
\nsolving constrained optimization problem with equality constraints  
\n(b) Solve the following problem by Lagrange multiplier method [5]  
\nOptimize 
$$
z = x_1^2 - 10x_1 + x_2^2 - 6x_2 + x_3^2 - 4x_3
$$
  
\nsubject to  $x_1 + x_2 + x_3 = 7$   
\n2. (a) Show that the Kuhn Tucker necessary conditions for the optimization problem  
\nMinimize  $f(x)$   
\nsubject to  $g_1(x) \le 0$ ,  $i = 1, 2, \dots, n$   
\nand  $h_j(x) = 0$ ,  $j = 1, 2, \dots, l$   
\nare also sufficient conditions if  $f(x)$  is convex and  $g_i(x)$ ,  $i = 1, 2, \dots, m$  are convex  
\nfunctions of x and  $h_j(x)$  are linear.  
\n(b) Use the Kuhn-Tucker necessary conditions to solve the following optimization problem:  
\nMaximize  $z = 2x_1 - x_1^2 + x_2$   
\nsubject to  $2x_1 + 3x_2 \le 6$ ,  $2x_1 + x_2 \le 4$  and  $3x_1 + 9x_2 = 16$   
\n3. (a)  
\nIf the iterative sequence  $\{x^{(k)}\}$  be defined as  
\n $x^{(k+1)} = x^{(k)} + \lambda^{(k)}d^{(k)}$   
\nwhere  $d^{(k)}$  is given by  
\n $d^{(k)} = -M_k \nabla f(x^{(k)})$   
\nand also, if  $\nabla f(x^{(k)}) \ne 0$  and  $M_k$  is positive definite, then show that the iterative procedure possesses  
\nprocedure possesses descent property.  
\n(b) Using steepest descent method, minimize  $f(x) = x_1^2 + x_2^2 + 2gx_1 + 2fx_2 + c$  starting [5]  
\nfrom the point  $\begin{pmatrix} \alpha \\ \beta \end{pmatrix}$ .

- 4. (a) Write down Wolfe's algorithm for solving quadratic programming problem. [4] (b) Use Beale's method for solving the quadratic programming problem: Maximize  $z = 4x_1 + 6x_2 - 2x_1^2 - 2x_1x_2 - 2x_2^2$ subject to  $x_1 + 2x_2 \le 2$  and  $x_1, x_2 \ge 0$ [5]
- 5. (a) Discuss Fibonacci search method for solving one-dimensional non-linear minimization problem. [4]
	- (b) Using Fibonacci search method

$$
\text{Maximize } f(x) = \begin{cases} \frac{2x+3}{6} & \text{for } x \le 3\\ -x+6 & \text{for } x > 3 \end{cases}
$$

in the interval  $[-1, 5]$  (consider 6 experimental points).

6. (a) Define Slater's, Karlin's and strict constraint qualifications. [3]

## (b) Show that

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- (i) Slater's constraint qualification and Karlin's constraint qualification are equivalent.
- (ii) The Strict constraint qualification implies Slater's constraint qualification and Kerlin's constraint qualification.
- 7. (a) State and prove Max-flow Min-cut theorem for a network flow problem. [4]

(b) Find the maximum flow in the graph with the following arcs and arc capacities, flow in each arc being non-negative. Arc  $(v_j, v_k)$  is denoted as  $(j, k)$ .  $v_a$  is the source and [5]

 $v_b$ , the sink.





# **M.A./M.Sc. Semester III Examination, 2019 (Old pattern under CDOE) Subject: Mathematics (Pure Stream)**

Time: 2 Hours Full Marks: 45

[5]

[3+3]

*The figures in the margin indicate full marks. Candidates are required to give their answers in their own words as far as practicable. [Notation and symbols have their usual meaning]*

#### **Paper: MPS305 (Operator Theory and Applications I)**

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Answer any **five** questions. Only **first five** answers will be evaluated.  $9 \times 5 = 45$ 

- 1. (a) Let *X* be a normed linear space with dual  $X^{\prime}$ . Define annihilator of a subset of *X* and  $X^{\prime}$ . [5] Show that the annihilator of a subset and the orthogonal complement of that subset coincide in a Hilbert space.
	- (b) Let  $T_1: X \to X$  and  $T_2: X \to X$  be two bounded linear transformations and  $T_1$  and  $T_2$  be its conjugates respectively where *X* is a normed linear space. Show that (i)  $(T_1+T_2)'$  $T_1' + T_2'$  and (ii)  $(T_1 T_2)' = T_2' T_1'.$ [4]
- 2. (a) Show that a surjective symmetric operator T:  $D_T \rightarrow X$  is self adjoint where  $D_T$  is a dense subspace of a complex Hilbert space *X.* [3]
	- (b) Let  $T \in B(X, Y)$  where X and Y are complex Hilbert spaces. Show that (i)  $(\overline{R(T)})^{\perp}$  $N(T^*)(ii) \overline{(R(T^*))}^{\perp} = N(T)(iii) ||T^*T|| = ||T||^2$ [6]
- 3. (a) Let  $E_1$  and  $E_2$  be two orthogonal projections on the closed subspaces  $M_1$  and  $M_2$  of a complex Hilbert space *X* respectively. Show that the following conditions are equivalent. (i) $E_1 \le E_2$ (ii)  $||E_1x|| \le ||E_2x||$  for all  $x \in X$  (iii)  $M_1 \subset M_2$  (iv)  $E_2E_1 = E_1$  (v)  $E_1E_2 =$  $E_1$ . [4]
	- (b) Let  $E_1$  and  $E_2$  be two orthogonal projections on the closed subspaces  $M_1$  and  $M_2$  of a complex Hilbert space *X* respectively. Prove that  $E_1 + E_2$  is an orthogonal projection on  $\overline{[M_1 \cup M_2]} = M_1 + M_2$  if and only if  $E_1$  is orthogonal to  $E_2$ . [5]
- 4. Let *X* be a complex Hilbertspace and let  $A, B \in B(X, X)$  and  $AB = BA$ . Show that  $A \geq 0, B \geq 0$  implies  $AB \geq 0$ . [9]
- 5. (a) If *A* and *B* are two compact operators defined on a normed linear space *X* into a normed linear space Y, then show that  $A + B$  is compact and  $\alpha A$  is compact where  $\alpha$  is a scalar. [4]
	- (b) Let  $H_1$  and  $H_2$  be two complex Hilbert spaces and let  $T: H_1 \rightarrow H_2$  be a compact linear operator. Prove that the adjoint operator  $T^*$  is compact. [5]
- 6. (a) Prove that the uniform limit of a sequence  ${A_n}$  of compact operators defined on a normed linear space *X* into a normed linear space *Y* is compact. [6]
	- (b) Let  $H_1$  and  $H_2$  be two normed linear spaces and let  $T: H_1 \rightarrow H_2$  be a compact linear operator. Show that range of *T* is separable. [3]
- 7. (a) Let *X* be a complex normed linear space and let  $A: X \to X$  be a compact linear operator and  $\lambda \neq 0$ . If  $y = \lambda x - Ax$  has a solution for arbitrary  $y \in X$ , then prove that the equation  $\lambda x - Ax = 0$  has only trivial solution  $x = 0$ . [6]
	- (b) Let *X* be a complex normed linear space and let  $A: X \rightarrow X$  be a compact linear operator and  $\lambda \neq 0$ . Then show that  $y = \lambda x - Ax$  is solvable for those and only those y belonging to  $a_{N(\lambda I - A')}$ . [3]